

Card-Based ZKP Protocol for Nurimisaki

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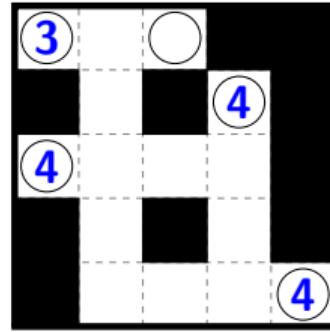
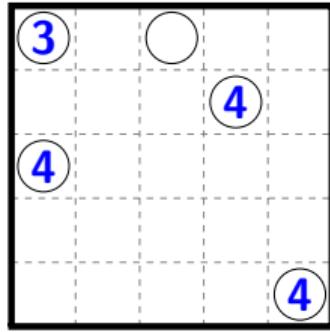
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November 17, 2022

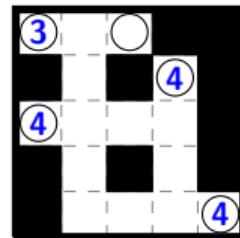
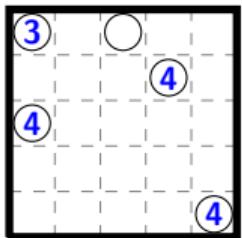


Nurimisaki



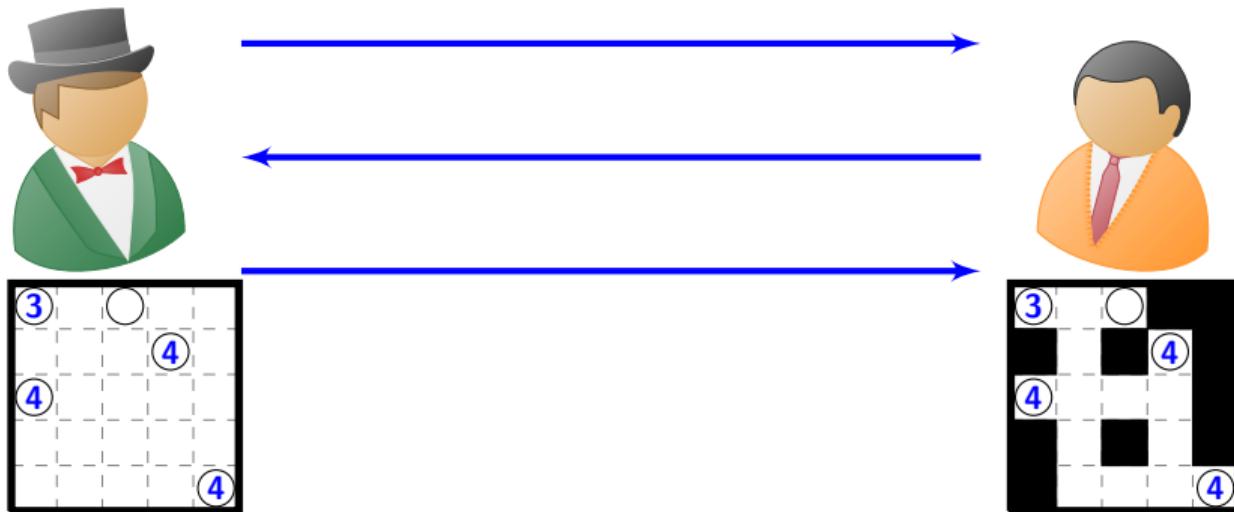
- ▶ ○ ⇔ Misaki: only one neighbour is white.
- ▶ **Number Rule:** Number of white cells in straight line from Misaki.
- ▶ **Square Rule:** No 2×2 square is composed of all blacks/whites.
- ▶ **Connectivity Rule:** White cells are connected.

Zero-Knowledge Proof (ZKP) Protocol



NP-complete

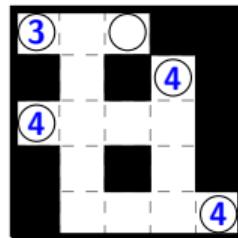
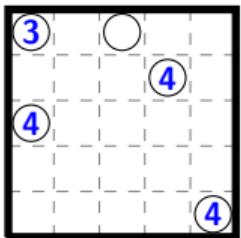
Zero-Knowledge Proof (ZKP) Protocol



Zero-Knowledge Proof (ZKP) Protocol



He has the solution.



ZKP properties

Completeness

If P knows the solution then it can convince V .

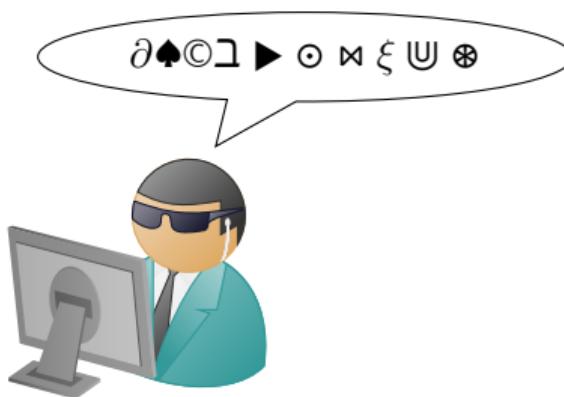
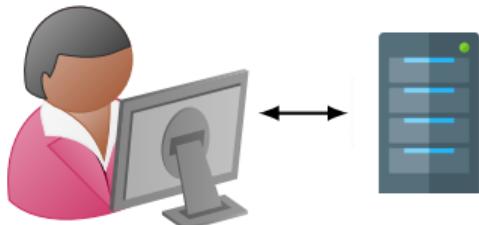
Extractability

If P does not have the solution then V can notice it.

Zero-Knowledge

V learns nothing about the solution.

ZKP for real...

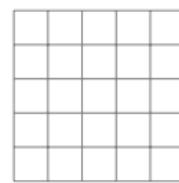


Encoding

0/black \longrightarrow 

1/white \longrightarrow 

dummy \longrightarrow 



G

0	1	0	1	1
0	0	1	0	1
1	1	0	1	0
1	0	1	0	0
0	1	1	0	1

S

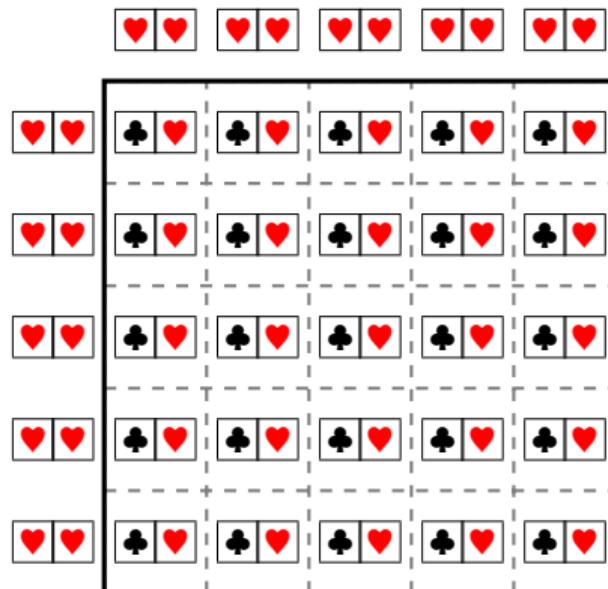
Shuffle



$$\left\langle \underbrace{\begin{array}{|c|} \hline ? \\ \hline \end{array}}_{p_1} \parallel \underbrace{\begin{array}{|c|} \hline ? \\ \hline \end{array}}_{p_2} \parallel \cdots \parallel \underbrace{\begin{array}{|c|} \hline ? \\ \hline \end{array}}_{p_m} \right\rangle \rightarrow \underbrace{\begin{array}{|c|} \hline ? \\ \hline \end{array}}_{p_{s+1}} \quad \underbrace{\begin{array}{|c|} \hline ? \\ \hline \end{array}}_{p_{s+2}} \quad \cdots \quad \underbrace{\begin{array}{|c|} \hline ? \\ \hline \end{array}}_{p_{s+m}}$$

Our protocol - Setup phase

P and V place black commitments on each cell + dummy around G



Roadmap

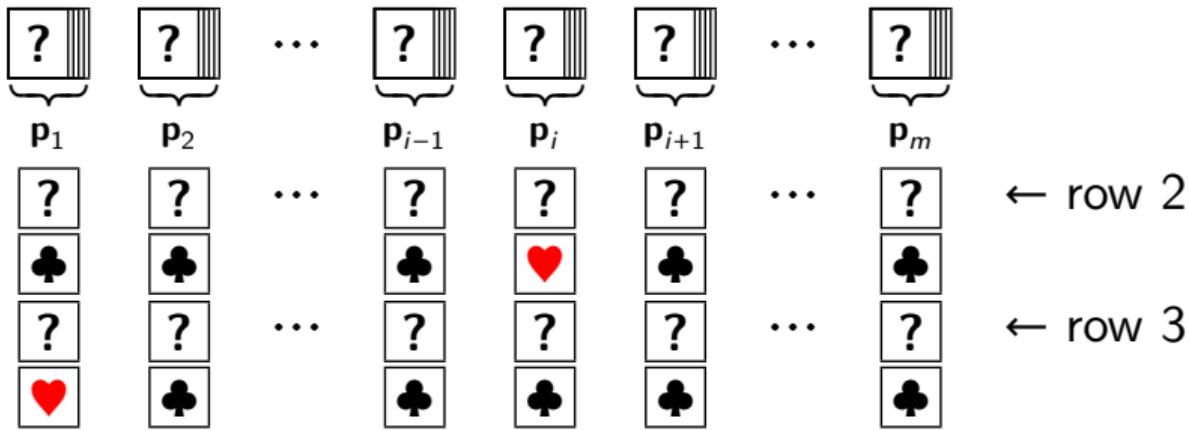
Connectivity

- ▶ Chosen Pile Protocol;
- ▶ 4-Neighbour Protocol;
- ▶ Forming the polyomino.

Verifications

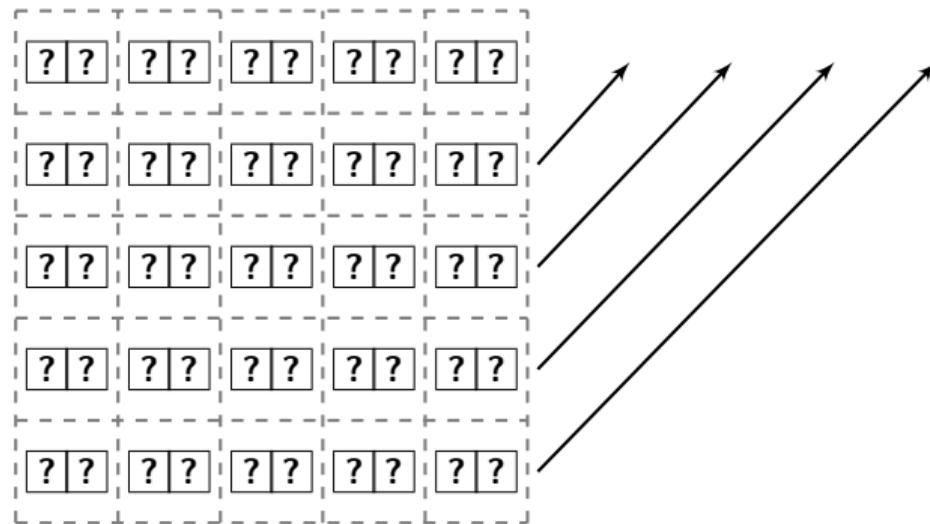
- ▶ No 2×2 square;
- ▶ Misaki;
- ▶ No circle \Leftrightarrow No Misaki.

Chosen Pile Protocol



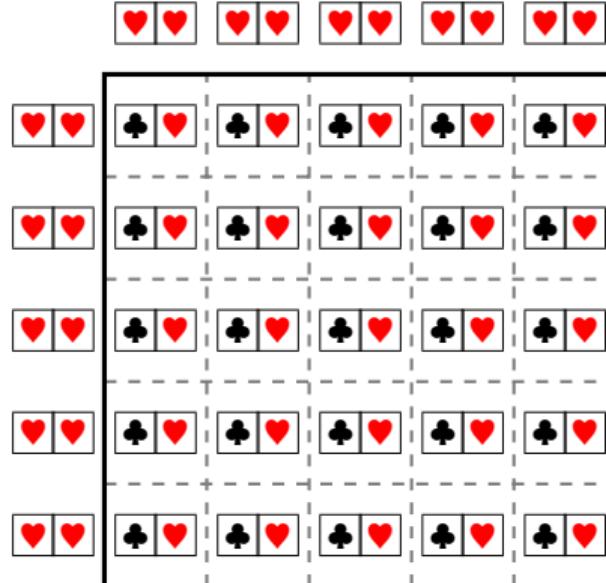
J-G. Dumas, P. Lafourcade, D. Miyahara, T. Mizuki, T. Sasaki, H. Sone,
Interactive Physical Zero-Knowledge Proof for Norinori. 166-177 2019 COCOON

Choose a Neighbour



? ? ? ? ? ? ? ? ? ?

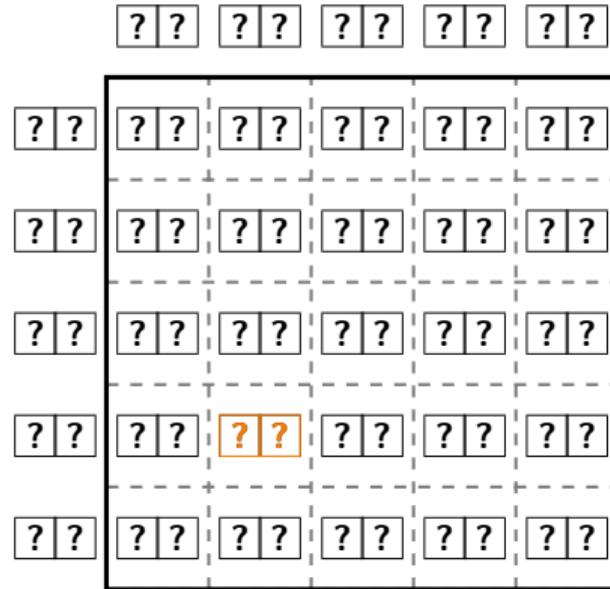
Forming White Polyomino



L. Robert, D. Miyahara, P. Lafourcade, T. Mizuki:

Card-Based ZKP for Connectivity: Applications to Nurikabe, Hitori, and Heyawake. New Gener. Comput. 40(1): 149-171 (2022)

Forming White Polyomino



P: Chosen Pile Protocol

V: Swap to white

L. Robert, D. Miyahara, P. Lafourcade, T. Mizuki:

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Forming White Polyomino



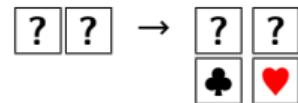
P: 4-Neighbour Protocol

L. Robert, D. Miyahara, P. Lafourcade, T. Mizuki:

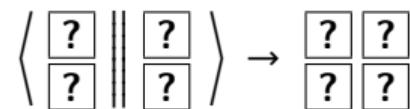
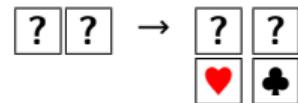
Card-Based ZKP for Connectivity: Applications to Nurikabe, Hitori, and Heyawake. New Gener. Comput. 40(1): 149-171 (2022)

Changed/Still Commitment

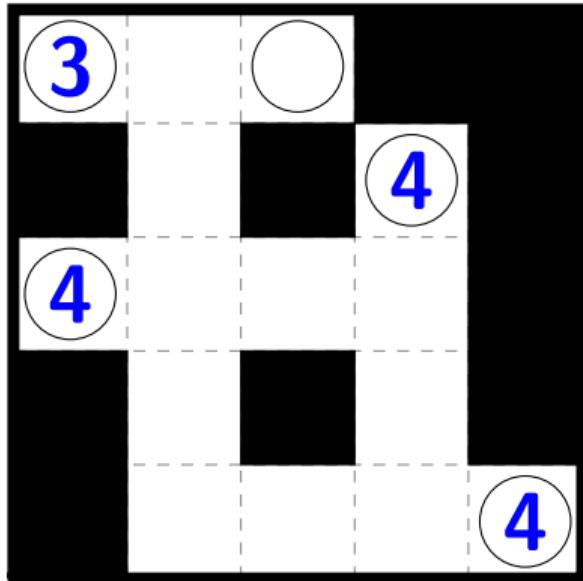
Change commitment:



Still commitment:



Verifications of other rules



V is convinced that white cells are connected.

V checks that other rules are satisfied:

- ▶ 2×2 square;
- ▶ Misaki cells have only one white neighbour;
- ▶ Number indicates the number of white cells in straight line;
- ▶ No Circle \Leftrightarrow No Misaki.

No 2×2 square



No 2×2 square



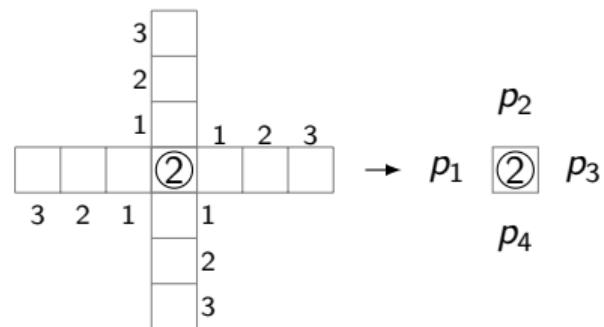
P: Chosen Pile Protocol

No 2×2 square



Misaki

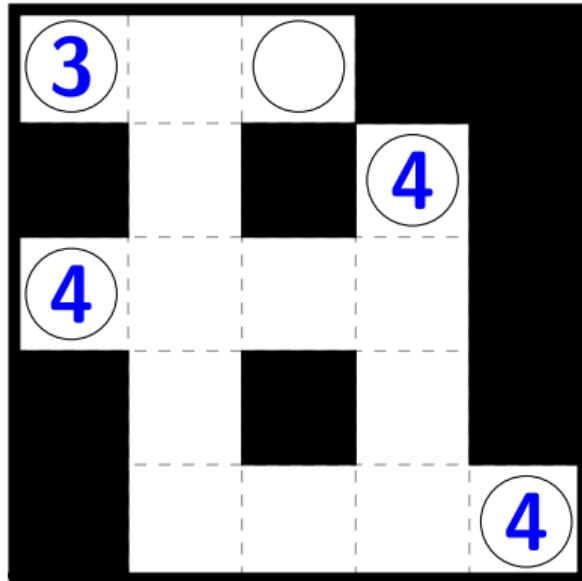
1. Only one neighbour is white;
2. Number indicates the number of white cells in straightline.



$$\left\langle \begin{array}{|c|} \hline p_1 \\ \hline 1 \\ \hline \end{array} \parallel \begin{array}{|c|} \hline p_2 \\ \hline 2 \\ \hline \end{array} \parallel \begin{array}{|c|} \hline p_3 \\ \hline 3 \\ \hline \end{array} \parallel \begin{array}{|c|} \hline p_4 \\ \hline 4 \\ \hline \end{array} \right\rangle \rightarrow \begin{array}{|c|} \hline ? \\ \hline ? \\ \hline ? \\ \hline ? \\ \hline \end{array} \begin{array}{|c|} \hline ? \\ \hline ? \\ \hline ? \\ \hline ? \\ \hline \end{array} \begin{array}{|c|} \hline ? \\ \hline ? \\ \hline ? \\ \hline ? \\ \hline \end{array} \begin{array}{|c|} \hline ? \\ \hline ? \\ \hline ? \\ \hline ? \\ \hline \end{array}$$

→ V: reveal 1st commitments

No Circle \Leftrightarrow No Misaki



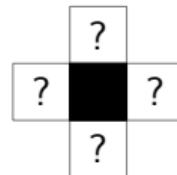
White Neighbours ≥ 2

⌚ Location and Number of white cells \subset solution.

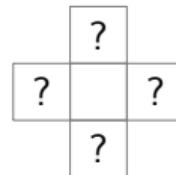
Idea: Sum

$$\blacksquare \longrightarrow \spadesuit \heartsuit \longrightarrow \spadesuit \heartsuit \heartsuit \heartsuit \heartsuit \heartsuit = 0$$

$$\square \longrightarrow \heartsuit \clubsuit \longrightarrow \heartsuit \heartsuit \heartsuit \heartsuit \heartsuit \clubsuit = 5$$



$$\sum \leq 4$$



$$\sum \leq 7$$

Security Proofs

Theorem (Completeness)

If P knows the solution of a Nurimisaki grid, then P can convince V .

→ We show that all the steps will not abort.

Theorem (Soundness)

If P does not provide a solution of the $p \times q$ Nurimisaki grid, then P is not able to convince V .

→ We show that at least one step will abort.

Theorem (Zero-knowledge))

V learns nothing about P 's solution of the given grid G .

→ Description of a simulator which does not have the solution but can swap cards during shuffles.

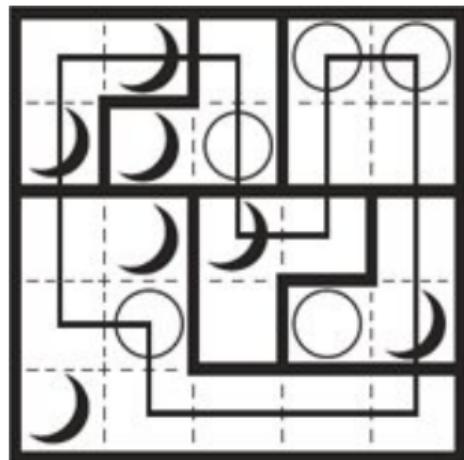
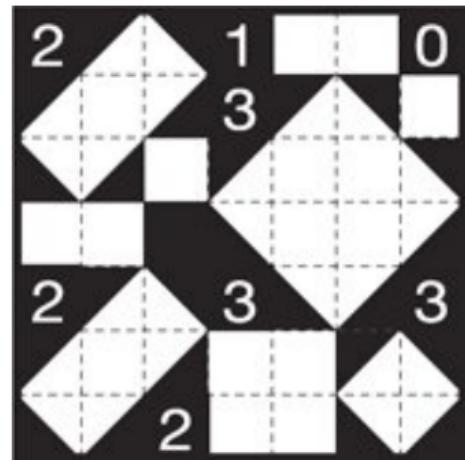
Conclusion

ZKP protocol: how to convince of having a secret without revealing it.

3 properties:

- ▶ Completeness;
- ▶ Soundness;
- ▶ Zero-Knowledge.

Future Works



Thank you for your attention, questions ?

