## Security Analysis of Electronic Exams

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CEA, 14th April 2016

## Exam



Filippo Galanti (Sora in Caserta 1852 - Buenos Aires 1953)

## Traditional Exam





Information technology for the assessment of knowledge and skills.













## Threats...









- Candidate cheating
- ▶ Bribed, corrupted or unfair examiners
- Outside attackers

## ...and their Mitigation

Most existing e-exam systems assume **trusted authorities** and focus on **student cheating**:

Exam centers

▶ Software solutions, e.g. ProctorU





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Exam centers

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#### Yet also the other threats are real:

- Atlanta Public Schools cheating scandal (2009)
- ► Turkish Public Personnel Selection Exam (2010)
- ▶ UK student visa tests fraud (2014)

## ...and their Mitigation

Most existing e-exam systems assume **trusted authorities** and focus on **student cheating**:

Exam centers

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- Atlanta Public Schools cheating scandal (2009)
- Turkish Public Personnel Selection Exam (2010)
- ▶ UK student visa tests fraud (2014)

So what about security of e-exams?

### Our Results

- Secrypt'14 Authentication Properties: Mark Authenticity, Answer Origin Authentication, Form Authorship, Form Authenticity.
  - Privacy Properties: Anonymous Marking, Question Indistinguishability, Anonymous Examiner, Mark Privacy, Mark Anonymity
- ISPEC'15 Individual Verifiability: Question Validity, Marking
  Correctness, Exam-Test Integrity, Exam-Test
  Markedness, Marking Integrity, Marking Notification
  Integrity
  Universal Verifiability: Eligibility (Registration),
  Marking Correctness Exam-Test Integrity, Exam-Test
  Markedness, Marking Integrity.
  - RV'15 How can we use previous results on real e-exam? Monitoring of reals e-exams.

### Plan

#### Introduction

### Security

Authentication Properties

**Privacy Properties** 

Huszti & Pethő's Protocol

Remark! Protocol

### Verifiability

Model

Grenoble Exam

Remark! Protocol

### Monitoring

Model

**Properties** 

Case Study: UJF E-exam

#### Conclusion

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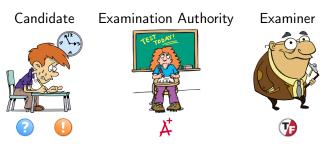
Properties

Case Study: UJF E-exam

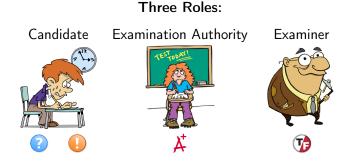
Conclusion

# E-exam: Players and Organization





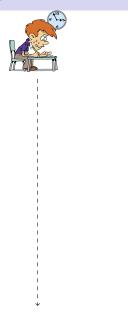
# E-exam: Players and Organization



#### Four Phases:

1. Registration 2. Examination 3. Marking 4. Notification

- ▶ **Processes** in the applied  $\pi$ -calculus [AF01]
- Annotated using events
- Authentication properties as correspondence between events
- Privacy properties as observational equivalence between instances
- ► Automatic verification using ProVerif [Bla01]













1. Registration







1. Registration Register register







1. Registration Register register

2. Examination

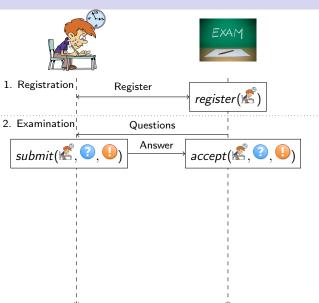






1. Registration Register register

2. Examination Questions

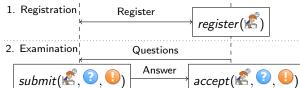




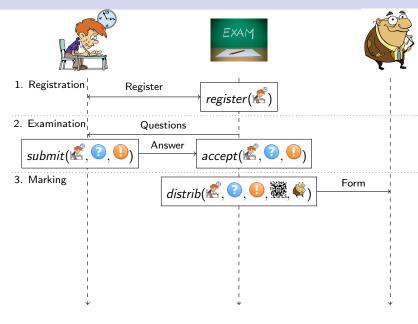


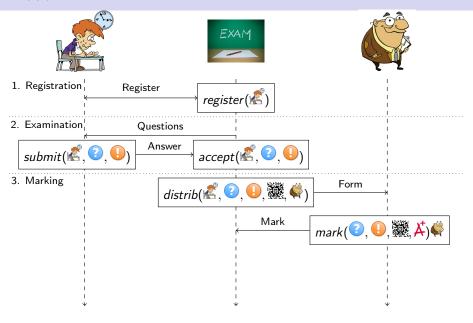


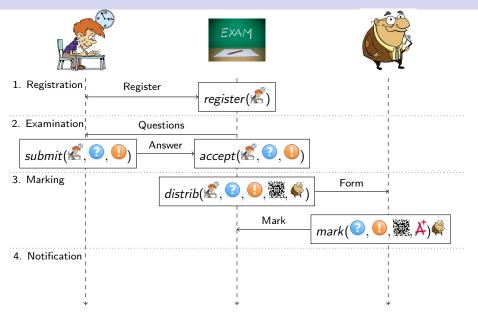


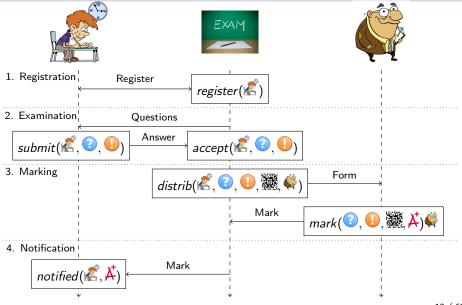


3. Marking









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Privacy Properties

Huszti & Pethő's Protocol

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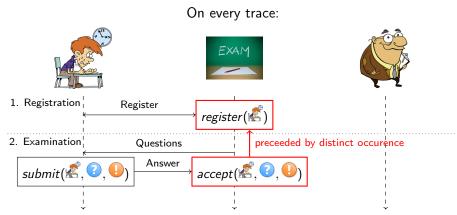
Properties

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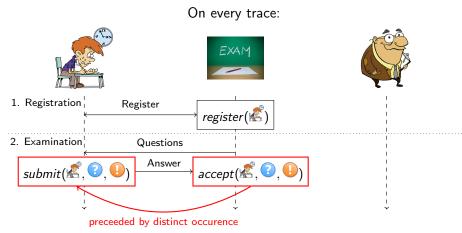
# Answer Origin Authentication

All collected answers originate from registered candidates, and only one answer per candidate is accepted.



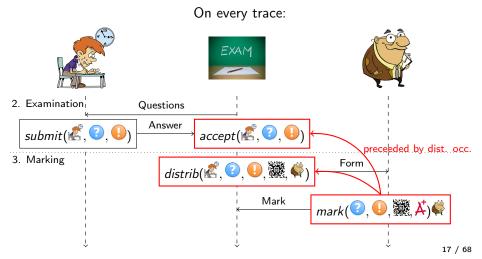
## Form Authorship

Answers are collected as submitted, i.e. without modification.



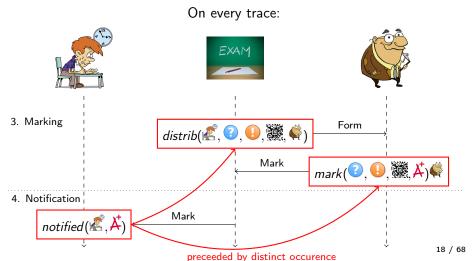
## Form Authenticity

Answers are marked as collected.



## Mark Authenticity

The candidate is notified with the mark associated to his answer.



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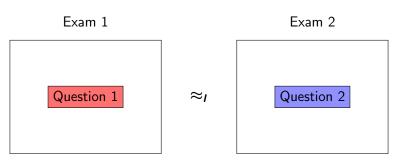
Conclusion

# Question Indistinguishability

No premature information about the questions is leaked.

#### **Definition:**

Observational equivalence of two instances up to the end of registration phase:

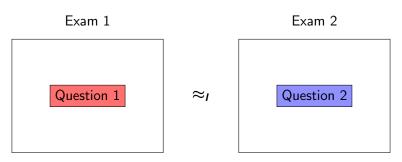


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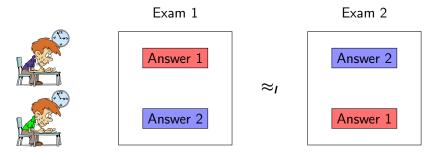
Can be considered with or without dishonest candidates.

# Anonymous Marking

An examiner cannot link an answer to a candidate.

#### **Definition:**

Up to the end of marking phase:

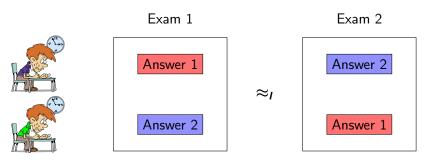


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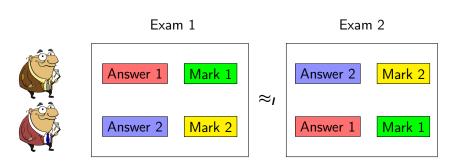


Can be considered with or without dishonest examiners and authorities.

# Anonymous Examiner

A candidate cannot know which examiner graded his copy.

#### **Definition:**

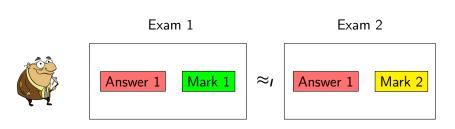


Can be considered with or without dishonest candidates.

# Mark Privacy

Marks are private.

#### **Definition:**

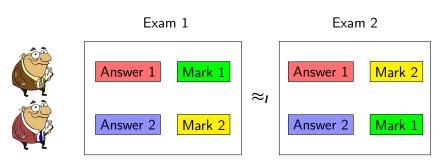


Can be considered with or without dishonest candidates, examiners and authorities.

# Mark Anonymity

Marks can be published, but may not be linked to candidates.

#### **Definition:**



Can be considered with or without dishonest candidates, examiners and authorities.

Implied by Mark Privacy.

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Remark! Protocol

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## Application: Huszti & Pethő's Protocol

## "A Secure Electronic Exam System" [HP10] using

- ElGamal Encryption
- a Reusable Anonymous Return Channel (RARC) [GJ03] for anonymous communication
- a network of servers providing a timed-release service using Shamir's Secret Sharing:

A subset of servers can combine their shares to **de-anonymize a candidate** after the exam

#### Goal: ensure

authentication and privacy

#### in presence of dishonest

- candidates
- examiners
- exam authorities

## Results

## Formal Verification with ProVerif [Bla01]:

Property	Result	Time
Answer Origin Authentication	×	< 1 s
Form Authorship	×	< 1 s
Form Authenticity	×	< 1 s
Mark Authenticity	×	< 1 s
Question Indistinguishability	×	< 1 s
Anonymous Marking	×	8 m 46 s
Anonymous Examiner	×	9 m 8 s
Mark Privacy	×	39 m 8 s
Mark Anonymity	×	1h 15 m 58 s

#### Main reason

#### Given its security definition, the RARC

- provides anonymity, but not necessarily secrecy
- does not necessarily provide integrity or authentication
- ▶ is only secure against passive attackers

Corrupted parties or active attackers can **break secrecy and anonymity**, as the following attack shows.

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## Application: Remark! Protocol

### A recent protocol [GLR14] using

- ElGamal encryption
- an exponentiation mixnet [HS11] to create pseudonyms based on the parties' public keys
  - ⇒ allows to encrypt and sign anonymously
- a public append-only bulletin board

#### Goal: ensure

- authentication and integrity
- privacy
- verifiability

### in presence of dishonest

- candidates
- examiners
- exam authorities

## Results

### Formal Verification with ProVerif:

Property	Result	Time
Answer Origin Authentication	✓	< 1 s
Form Authorship	<b>√</b>	< 1 s
Form Authenticity	$\checkmark^1$	< 1 s
Mark Authenticity	✓	< 1 s
Question Indistinguishability	<b>√</b>	< 1 s
Anonymous Marking	✓	2 s
Anonymous Examiner	✓	1 s
Mark Privacy	<b>√</b>	3 m 32 s
Mark Anonymity	✓	_2

<sup>&</sup>lt;sup>1</sup>after fix

<sup>&</sup>lt;sup>2</sup>implied by Mark Privacy

```
Huszti & Pethő's Protocol
Verifiability
```

```
Huszti & Pethő's Protocol
Verifiability
   Model
```

#### Exam model

### Very abstract model:

- Four sets:
  - $\{\mathcal{K}\}$ : candidate identities, subset  $\{\mathcal{K}\}_r$  registered candidates
  - $\{0\}$ : questions, subset  $\{0\}_g$  correct questions
  - ▶ {**!**}: answers
  - ► {♣}: marks
- ► Three relations:
  - ► Accepted  $\subseteq \{\cancel{\&}\} \times (\{^{2}\} \times \{^{1}\})$
  - ▶ Marked  $\subseteq \{ \checkmark \end{cases} \times (\{ \circlearrowleft \} \times \{ \circlearrowleft \}) \times \{ \checkmark \}$
  - Assigned  $\subseteq \{ \stackrel{\longleftarrow}{\mathbb{A}} \} \times \{ \stackrel{\longleftarrow}{\mathbb{A}} \}$
- ▶ A function Correct :  $(\{ {}^{\bigodot} \} \times \{ {}^{\bigodot} \}) \rightarrow \{ \stackrel{\longleftarrow}{\not} \}$
- ► An exam protocol is X-verifiable, if we have a sound and complete test for X.

# Defining Individual Verifiability

#### Each candidate knows

- her identity K<sup>6</sup>,
- ▶ question ②,
- ▶ answer ⊕,
- ► mark Å,
- ▶ and a log 🖼.

#### Properties:

The candidate can verify that...

Question Validity: ...she received questions generated by the question committee

$$QV_{IV}(\cancel{\&}, \cancel{Q}, \cancel{Q}, \cancel{A}, \cancel{\boxtimes}) \Leftrightarrow (\cancel{Q} \in \{\cancel{Q}\}_g)$$

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$$\mathsf{QV}_{\mathsf{IV}}(\overset{\bullet}{\mathbb{K}}, \overset{\bullet}{\mathbf{Q}}, \overset{\bullet}{\mathbf{Q}}, \overset{\bullet}{\mathbb{K}}, \overset{\bullet}{\mathbb{K}}) \Leftrightarrow (\overset{\bullet}{\mathbf{Q}} \in \{\overset{\bullet}{\mathbf{Q}}\}_g)$$
sound & complete

# Defining Individual Verifiability Cont'd

The candidate can verify that...

Marking Correctness: ...the mark attributed to her answer is correct.

$$\mathtt{MC}_{\mathtt{IV}}(\overset{\bullet}{\mathbb{A}}, \overset{\bullet}{\mathbf{O}}, \overset{\bullet}{\mathbf{O}}, \overset{\bullet}{\mathbf{A}}, \overset{\bullet}{\mathbf{D}}) \Leftrightarrow (\mathtt{Correct}(\overset{\bullet}{\mathbf{O}}, \overset{\bullet}{\mathbf{O}}) = \overset{\bullet}{\mathsf{A}})$$

► Exam-Test Integrity: ...her answer was accepted and marked as submitted.

$$\mathsf{ETI}_{\mathsf{IV}}(\mathbf{A}, \mathbf{O}, \mathbf{O}, \mathbf{A}, \mathbf{B}) \Leftrightarrow ((\mathbf{A}, (\mathbf{O}, \mathbf{O})) \in \mathsf{Accepted} \land \exists m' : (\mathbf{A}, (\mathbf{O}, \mathbf{O}), m') \in \mathsf{Marked})$$

► Exam-Test Markedness: ...her answer was marked.

$$\mathsf{ETM}_{\mathsf{IV}}(\mathbf{k}, \mathbf{2}, \mathbf{0}, \mathbf{k}, \mathbf{m}) \Leftrightarrow (\exists m' : (\mathbf{k}, (\mathbf{2}, \mathbf{0}), m') \in \mathsf{Marked}))$$

# Defining Individual Verifiability Cont'd

The candidate can verify that...

► Marking Integrity: ...her registered mark is the one assigned by the examiner

$$\mathtt{MI}_{\mathtt{IV}}(\nwarrow, \circlearrowleft, \overset{\bullet}{\bullet}, \overset{\bullet}{\longleftarrow}) \Leftrightarrow \exists m' : ((\nwarrow, (\overset{\bullet}{\bullet}, \overset{\bullet}{\bullet}), m') \in \mathtt{Marked} \land (\nwarrow, m') \in \mathtt{Assigned})$$

Marking Notification Integrity: ...she received the assigned mark

$$\mathtt{MNI}_{\mathtt{IV}}(\mathscr{A}, ?), \mathsf{U}, \mathsf{A}, \mathsf{w}) \Leftrightarrow (\mathscr{A}, \mathsf{A}) \in \mathtt{Assigned}$$

# Universal Verifiability

An **outside auditor** only has access to some evidence ...

The auditor can verify that...

### Properties:

Registration: ...all the accepted answers were submitted by registered candidates.

$$R_{UV}(igotimes_i) \Leftrightarrow \{igotimes_i^{\infty}\}_r \supseteq \langle i:(i,x) \in Accepted \rangle$$

► Marking Correctness: ...all the marks were calculated correctly.

$$ext{MC}_{ ext{UV}}(oxedowndots i) \Leftrightarrow orall (i, x, m) \in ext{Marked}, ext{Correct}(x) = m$$

## Universal Verifiability Cont'd

The auditor can verify that...

Exam-Test Integrity: ...all and only accepted test answers were marked.

$$\mathtt{ETI}_{\mathtt{UV}}(\stackrel{\square}{\mathtt{los}}) \Leftrightarrow \mathtt{Accepted} = \langle (i,x) : (i,x,m) \in \mathtt{Marked} \rangle$$

Exam-Test Markedness: ...all accepted test answers were marked.

$$\mathsf{ETM}_{\mathsf{UV}}(\stackrel{\mathsf{loc}}{=}) \Leftrightarrow \mathsf{Accepted} \subseteq \langle (i,x) : (i,x,m) \in \mathsf{Marked} \rangle$$

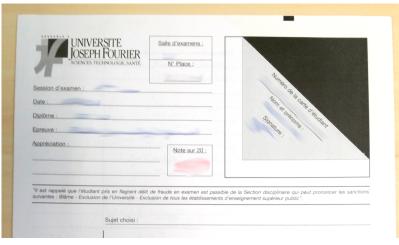
Marking Integrity: ...all and only the marks assigned to test answers were registered.

$$\mathtt{MI}_{\mathtt{UV}}(\overset{\square}{\mathtt{Dd}}) \Leftrightarrow \mathtt{Assigned} = \langle (i, m) : (i, x, m) \in \mathtt{Marked} \rangle$$

```
Huszti & Pethő's Protocol
Verifiability
   Grenoble Exam
```

## Case Study I: Grenoble Exam

- Paper-based exam system at the University Joseph Fourier
- ► Goal: Privacy (Anonymous Marking)
- ► **Special exam paper** with corner that is folded and glued:



## Case Study I: Grenoble Exam

- Paper-based exam system at the University Joseph Fourier
- ► Goal: Privacy (Anonymous Marking)
- ▶ Special exam paper with corner that is folded and glued:



### Grenoble Exam: Results

### Individual Verifiability:

- Input: the candidate's values
- Assumptions: Correct is published after the exam, and candidates can consult their copies
- Verification using ProVerif:

Property	Sound	Complete
Question Validity	× (EA)	<b>√</b>
Test Answer Integrity	× (EA, E)	<b>√</b>
Test Answer Markedness	× (E)	<b>√</b>
Marking Correctness	✓	<b>√</b>
Mark Integrity	× (EA, E)	<b>√</b>
Mark Notification Integrity	× (EA)	<b>√</b>

No guarantee that the records are correct.

## Grenoble Exam: Results Cont'd

### Universal Verifiability:

- Assumption: the auditor gets access to the EA's and Es' records and the function Correct.
- Verification using ProVerif:

Property	Sound	Complete
Registration	× (EA)	<b>√</b>
Exam-Test Integrity	× (EA, E)	<b>√</b>
Exam-Test Markedness	× (EA, E)	<b>√</b>
Marking Correctness	× (E)	<b>√</b>
Mark Integrity	× (EA, E)	<b>√</b>

▶ No guarantee that the records are correct, EA and E can make up fake records as long as they are coherent.

```
Huszti & Pethő's Protocol
Verifiability
   Remark! Protocol
```

## Remark!: Results

### Individual Verifiability:

- ► Input: the candidate's values and the messages on the bulletin board
- Assumption: Correct is published after the exam
- Verification using ProVerif:

Property	Sound	Complete
Question Validity	× (EA)	✓
Test Answer Integrity	<b>√</b>	<b>√</b>
Test Answer Markedness	✓	<b>√</b>
Marking Correctness	× (EA)	<b>√</b>
Mark Integrity	✓	<b>√</b>
Mark Notification Integrity	<b>√</b>	<b>√</b>

## Remark!: Results Cont'd

### Universal Verifiability:

- ► Input: the messages on the bulletin board, the function Correct, as well as additional data from the EA
- Verification using ProVerif:

Property	Sound	Complete
Registration	✓	✓
Exam-Test Integrity	✓	✓
Exam-Test Markedness	✓	✓
Marking Correctness	× (EA)	✓
Mark Integrity	<b>√</b>	✓

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Huszti & Pethő's Protocol
Monitoring
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```

```
Huszti & Pethő's Protocol
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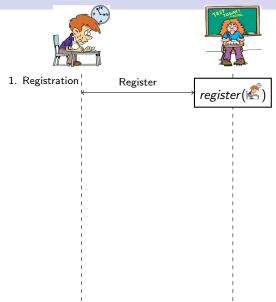








## 1. Registration





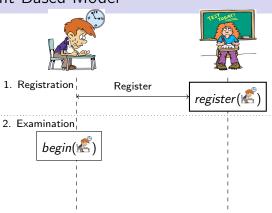




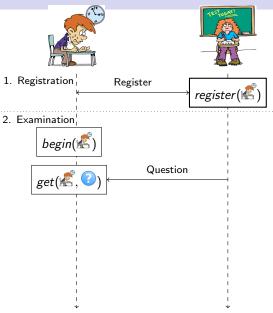


1. Registration Register register

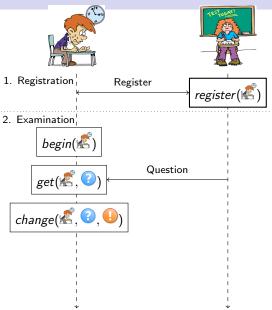
2. Examination



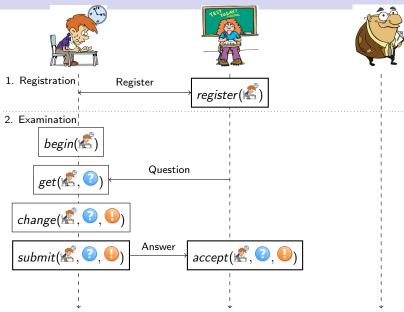


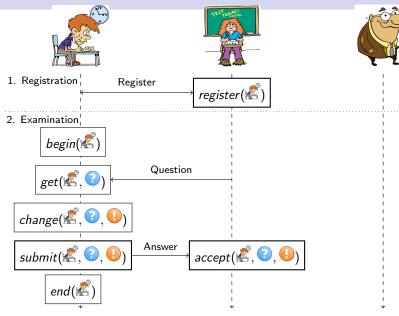














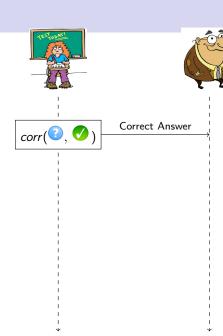




3. Marking |

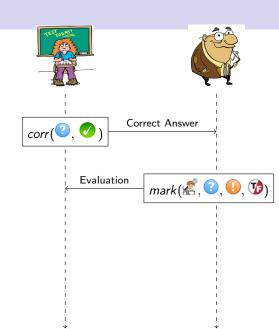


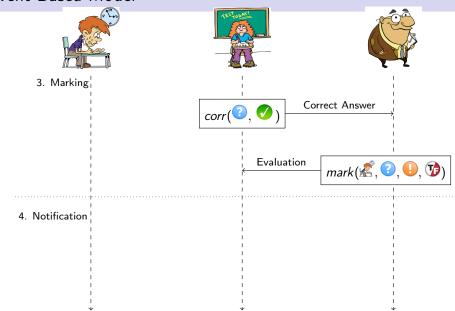
3. Marking 1

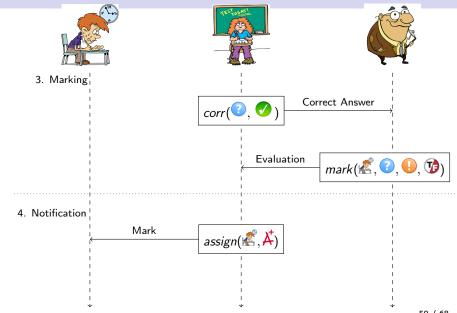




3. Marking







#### Plan

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Huszti & Pethő's Protocol
Monitoring
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# Quantified Event Automata (QEAs)

- Properties expressed as QEAs: event automaton with quantified variables.
- An event automaton is a finite-state machine with transitions labeled by parametric events.
- Transitions may include guards and assignments.
- ▶ We extend the initial definition of QEAs by:
  - 1. variable declaration and **initialization** before reading the trace
  - 2. **global variable** shared among all event automaton instances.
  - event(parameters) [guard] assignment

## Candidate Eligibility

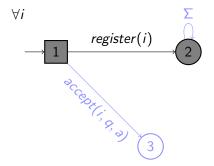
No answer is accepted from an unregistered candidate

$$\Sigma = \{register(i), accept(i, q, a)\}$$
 $\forall i$ 
 $register(i)$ 
 $2$ 

# Candidate Eligibility

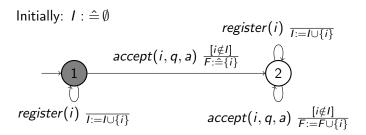
No answer is accepted from an unregistered candidate

$$\Sigma = \{ register(i), accept(i, q, a) \}$$



# Candidate Eligibility with Auditing

All candidates that violates the requirement are collected in a set F.



## Properties

Candidate Registration: an unregistered candidate tried to take the exam.

## **Properties**

Candidate Registration: an unregistered candidate tried to take the exam.

#### Answer Authentication:

- an unsubmitted answer was considered as accepted; or
- ▶ more than one answer were accepted from a candidate.

### **Properties**

Candidate Registration: an unregistered candidate tried to take the exam.

#### **Answer Authentication:**

- an unsubmitted answer was considered as accepted; or
- more than one answer were accepted from a candidate.

#### **Questions Ordering:**

► a candidate got a question before validating the previous ones.

Exam Availability: an answer was accepted outside exam time.

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#### Exam Availability with Flexibility:

 supports different duration and starting time between candidates.

Exam Availability: an answer was accepted outside exam time.

#### Exam Availability with Flexibility:

supports different duration and starting time between candidates.

Marking Correctness: an answer was marked in a wrong way.

Exam Availability: an answer was accepted outside exam time.

#### Exam Availability with Flexibility:

supports different duration and starting time between candidates.

Marking Correctness: an answer was marked in a wrong way.

#### Mark Integrity:

- an accepted answer was not marked; or
- a candidate was not assigned the corresponding mark.

#### Plan

```
Huszti & Pethő's Protocol
Monitoring
   Case Study: UJF E-exam
```

# E-exam at Université Joseph Fourier (UJF)

### Registration:

- 2 weeks before the exam.
- Using login/password.

# E-exam at Université Joseph Fourier (UJF)



#### Examination in a supervised room

Authentication and answers questions as follows:

- In a fixed order.
- ▶ Once validates the current question, he gets the next one.
- He can change the answer unlimited times before validating.
- Once he validates, then he cannot go back and change any of the validated answers.

# E-exam at Université Joseph Fourier (UJF)

#### Marking:

- ► For each question, the professor specifies the correct answer(s).
- ► For each question, all the answers provided by the candidates are collected.
- ▶ Each answer is evaluated by an examiner to 0 or 1.
- ► The mark for each candidate is calculated as the summation of all the scores attributed to his answers.

#### Notification:

- The marks are notified to the candidates.
- A candidate can consult his submission and check the marking.

## **Analysis**

Verification of two real e-exam executions using MarQ tool [RCR15].

From the logs: register(i), change(i, q, a), submit(i, q, a), accept(i, q, a).

### 4 Properties

- Candidate Registration
- Candidate Eligibility
- Answer Authentication
- Exam Availability

- ► Answer Authentication \*:
  - All accepted answers are submitted by candidates.
  - ▶ Allow the acceptance of the same answer again.
  - But, still forbids the acceptance of a different answer.

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- Answer Editing: A candidate cannot change an answer after validation it.

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- Question Ordering \*: A candidate cannot changes the answer to a future question before validating the current question.

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- Answer Editing: A candidate cannot change an answer after validation it.
- Question Ordering \*: A candidate cannot changes the answer to a future question before validating the current question.
- ► Acceptance Order: A candidate has to validate the questions in order, but he can skip some questions.

### Results: Exam 1

#### 233 students, 40875 events

Property	Result	Time (ms)
Candidate Registration	<b>√</b>	538
Candidate Eligibility	✓	517
Answer Authentication	×	310
Exam Availability	✓	518
Answer Authentication *	✓	742
Answer Authentication Reporting	×[1]	654
Answer Editing	✓	641
Question Ordering *	×	757
Acceptance Order	<b>√</b>	697

### Results: Exam 2

### 90 students, 4641 events

Property	Result	Time (ms)
Candidate Registration	✓	230
Candidate Eligibility	✓	214
Answer Authentication	✓	275
Exam Availability	×[1]	237
Answer Authentication *	✓	223
Answer Authentication Reporting	✓	265
Answer Editing	×	218
Question Ordering *	×	389
Acceptance Order	<b>√</b>	294

#### Plan

```
Huszti & Pethő's Protocol
```

#### Conclusion

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- Formal model for security and verifiability
- Security Analysis of 2 e-exams and one "real" exam
- Trust parties are required for verifiability
- Monitoring analysis of 2 real e-exams at UJF using MarQ tool
- Discovering some misbehaviours and flaws

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Designing secure protocols is difficult

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Designing secure protocols is difficult

Use formal methods!

#### Future Work

- ▶ Analyze more existing e-exams from other universities.
- Perform on-line verification with our monitors during live e-exams.
- ► Study more expressive and quantitative properties that can detect colluded students through similar answer patterns.
- Automatic transformation from verifiability to monitors. First try using a combination of model checking and monitoring.

# Thank you for your attention!

Questions?

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